
CSCE 4114 (Real Time) Operating Systems

David Andrews

dandrews@uark.edu



Operating Systems

- Originally developed to ease sharing of resources between users and foster portability
 - Early programming involved developing program specifically for a machine
 - Programs had to be re-written for each new machine.
- An OS is a "Virtual Machine"
 - Machine capabilities accessed through API's
 - User's code to API not machine specific registers, protocols, addresses, etc.
 - Specific implementations of API's provided through libraries
 - Libraries linked into source code



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Real Time Operating Systems

- Real-time OS: a multitasking[/multithreading] operating system for executing real-time applications.
 - Use specialized scheduling algorithms to deliver deterministic behavior.
 - Latency Considerations instead of throughput drives design. Sometimes miss-interpreted as "fast".
 - Typically modeled as event-driven: responds to change's in external environment such as input sensors
 - Event-driven system switches between tasks based on their priorities or external events while time-sharing operating systems switch tasks based on clock interrupts.



Embedded Operating Systems

- Embedded OS: Designed to operate on small machines like PDAs with less autonomy. They are able to operate with a limited number of resources. They are very compact and extremely efficient by design.
 - Small footprints
 - Scaled back capabilities
 - Virtual Memory Support



Operating System Services

- Program Management
 - Scheduling
- Timer Services
 - Date/Time
 - Watchdog Timers
- Synchronization/Communications
- File Services
- Networking
- Security



Scheduling

Task State

- 1 - has CPU

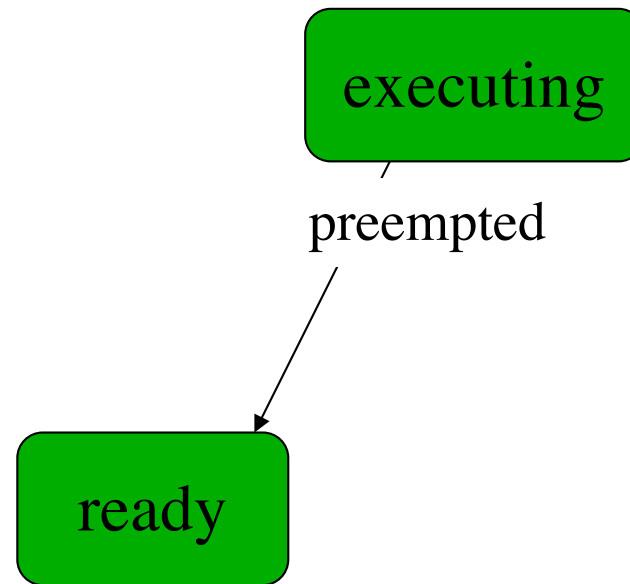
executing



Scheduling

Task State

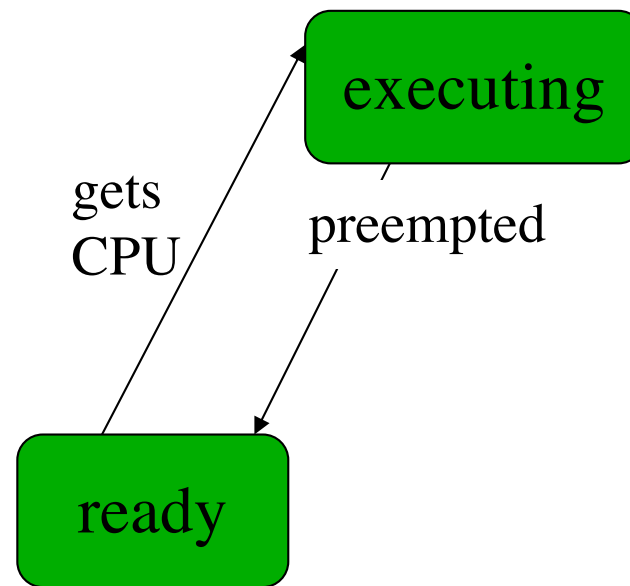
- 1 - has CPU
- No Longer Has CPU: Why ?
- Gets Preempted



Scheduling

Task State

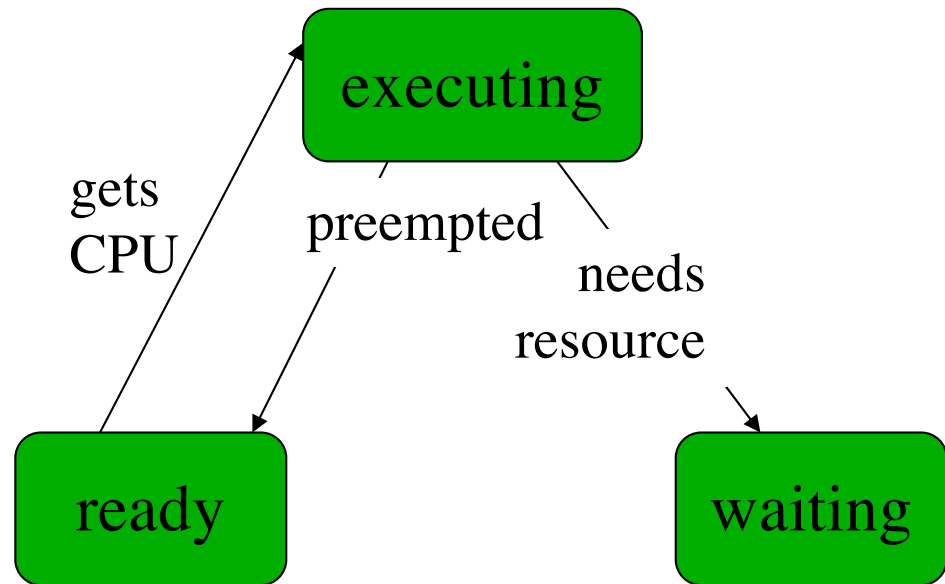
- 1 - has CPU
- No Longer Has CPU: Why ?
- Gets Preempted
- Still can run if possible



Scheduling

Task State

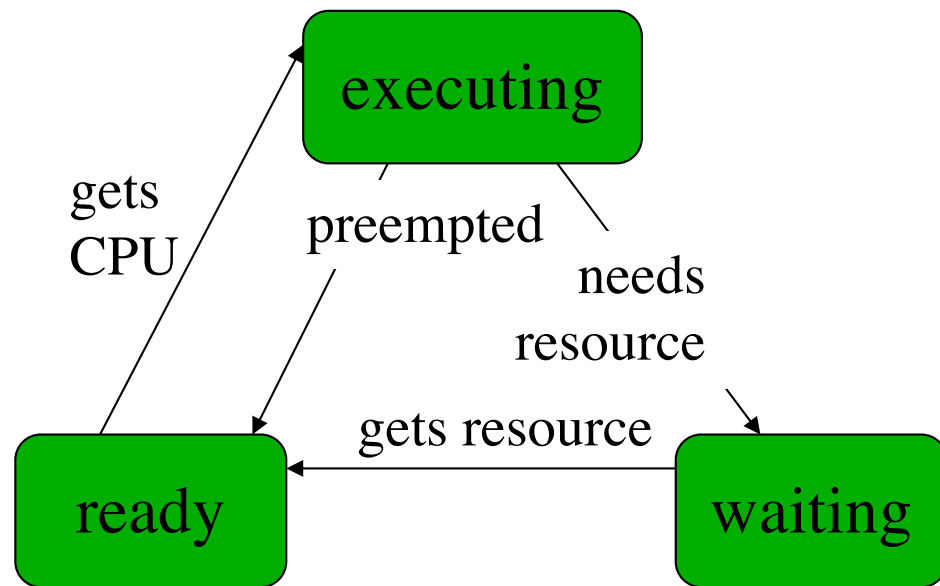
- 1 - has CPU
- No Longer Has CPU: Why ?
- Gets Preempted
- Needs to Wait on some resource
 - semaphore
 - I/O



Scheduling

Task State

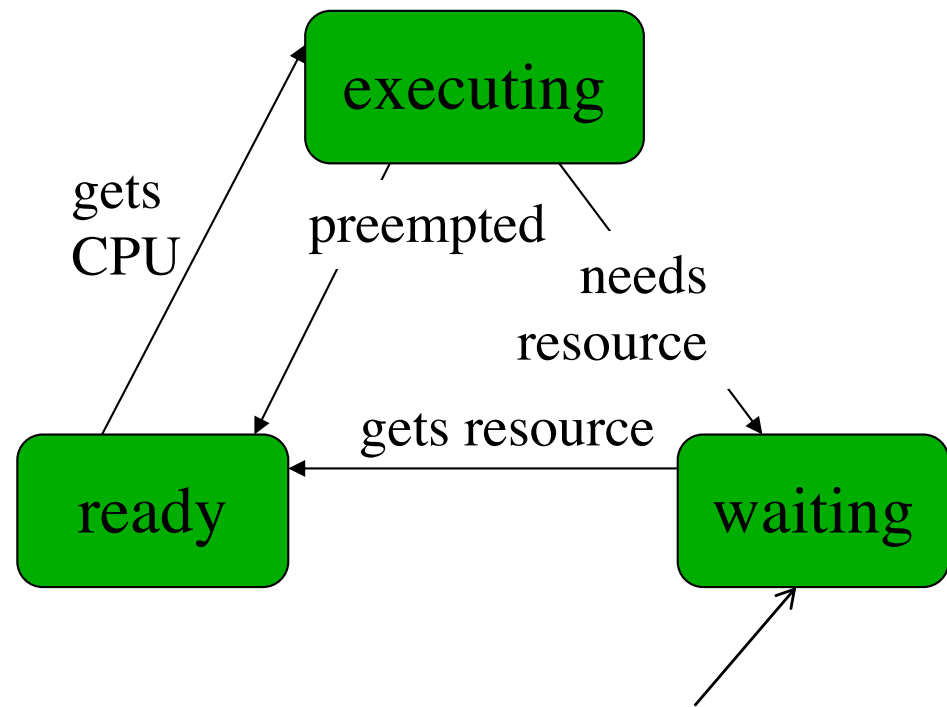
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- Gets Data and can now run



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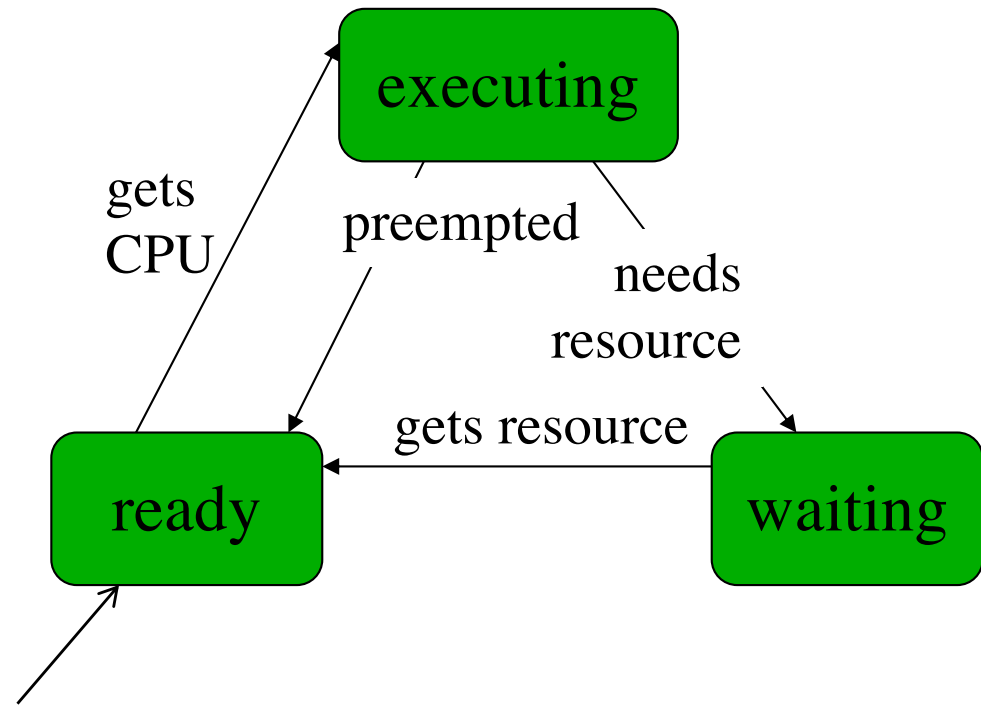
Called suspend or blocked Queue



Scheduling

Task State

- 1 - has CPU
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- Needs to Wait on some resource
 - semaphore
 - I/O
- Gets Data and can now run

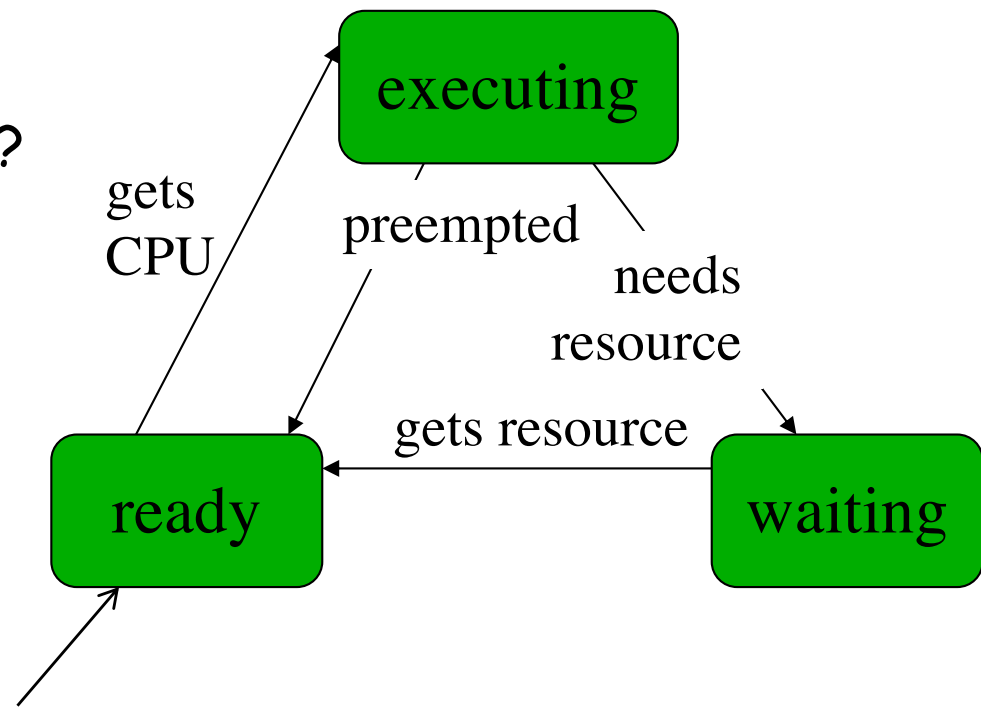


Called Ready to Run (R2R)
Or scheduler Queue



Scheduling

- If multiple threads in R2R Queue:
 - How does OS make scheduling decision?

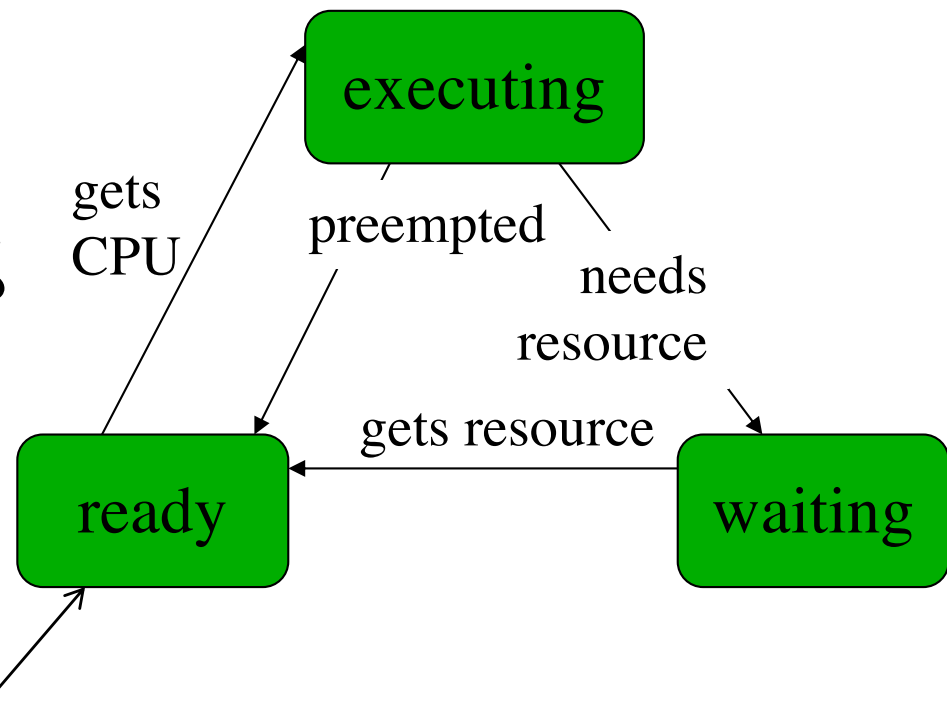


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Scheduling

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 - When does OS make scheduling decision?



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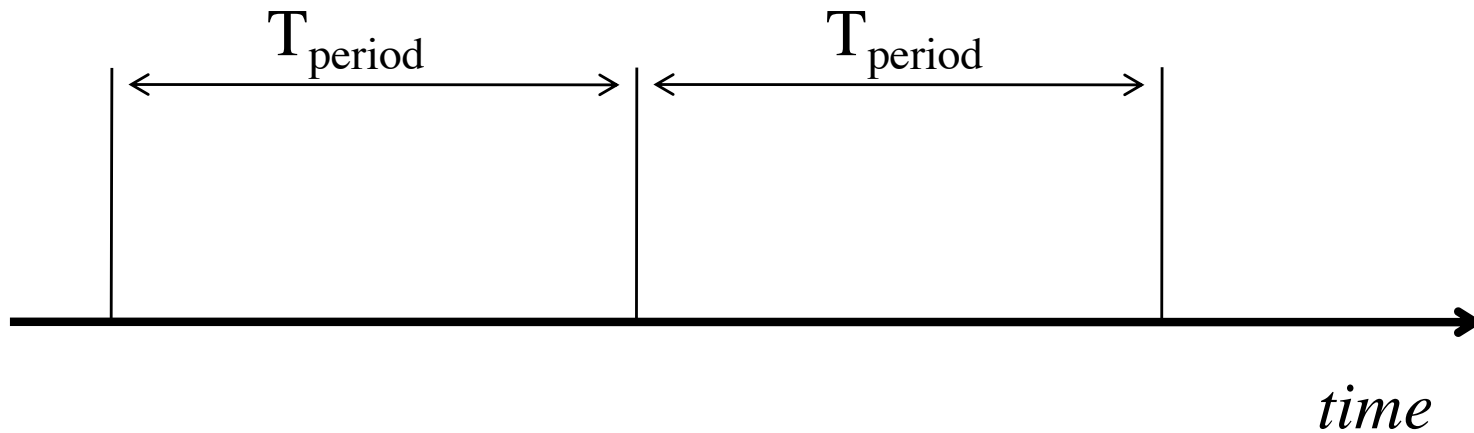


RT- Scheduling Algorithms

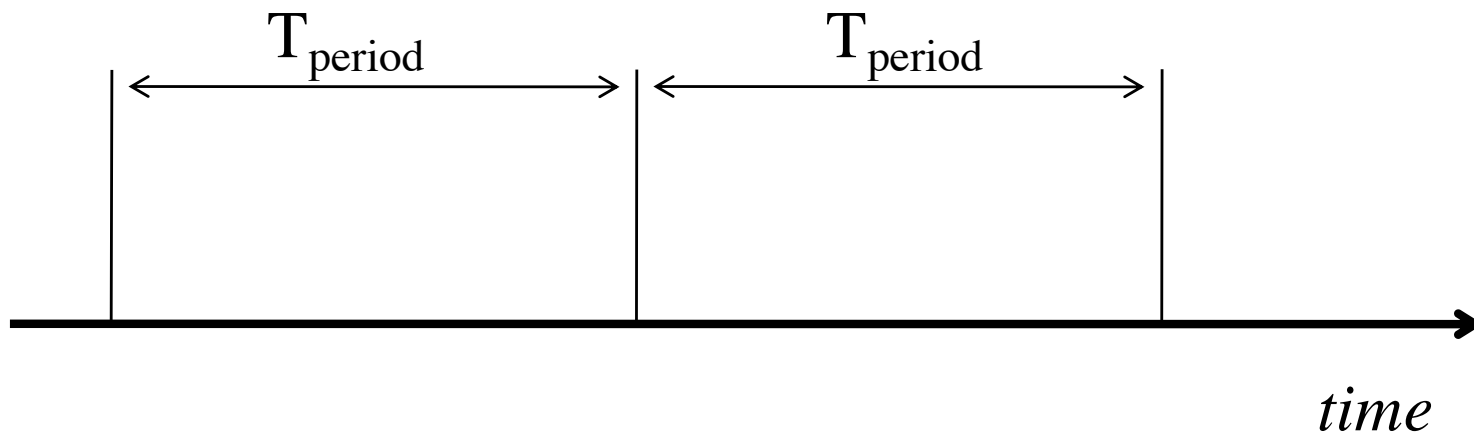
- Schedule multiple threads/tasks on shared resource(s) such that they all meet their deadlines.....
- Need to Know...
 - Execution time of each task
 - Study of Worst Case Execution Time
 - Deadline of each task
 - When all must be completed
 - When can task begin to execute
 - Periodic is simplest (aperiodic much more difficult)



A little theory....



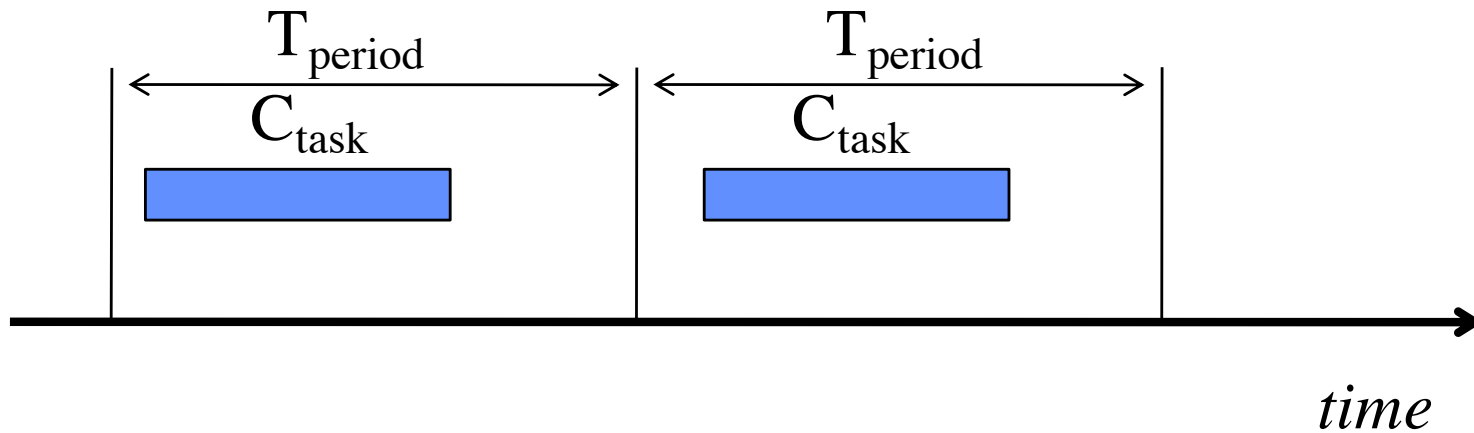
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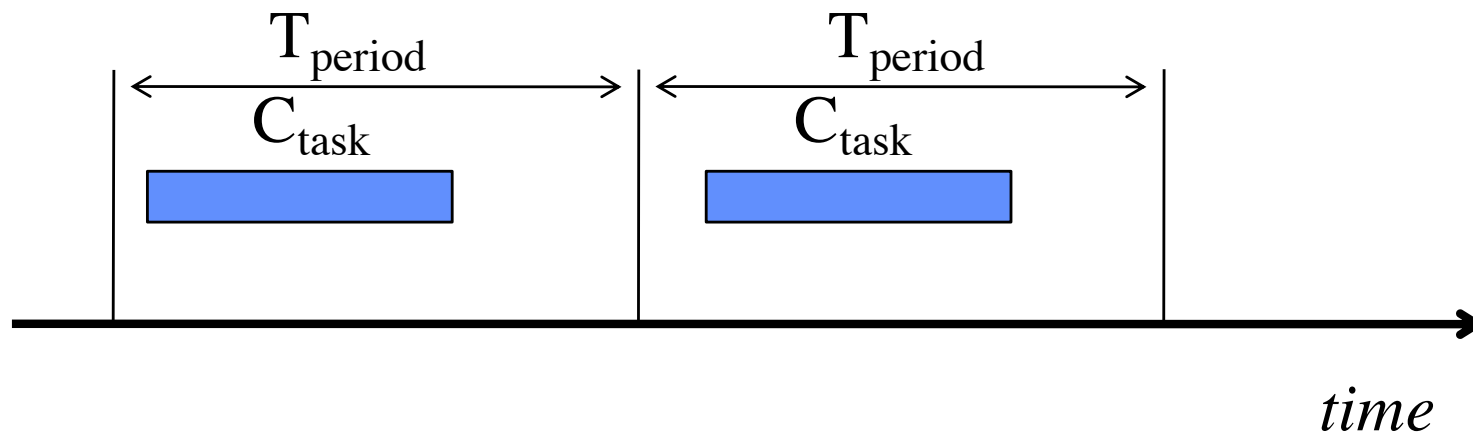
Task will “start” at the beginning of it’s period



A little theory....



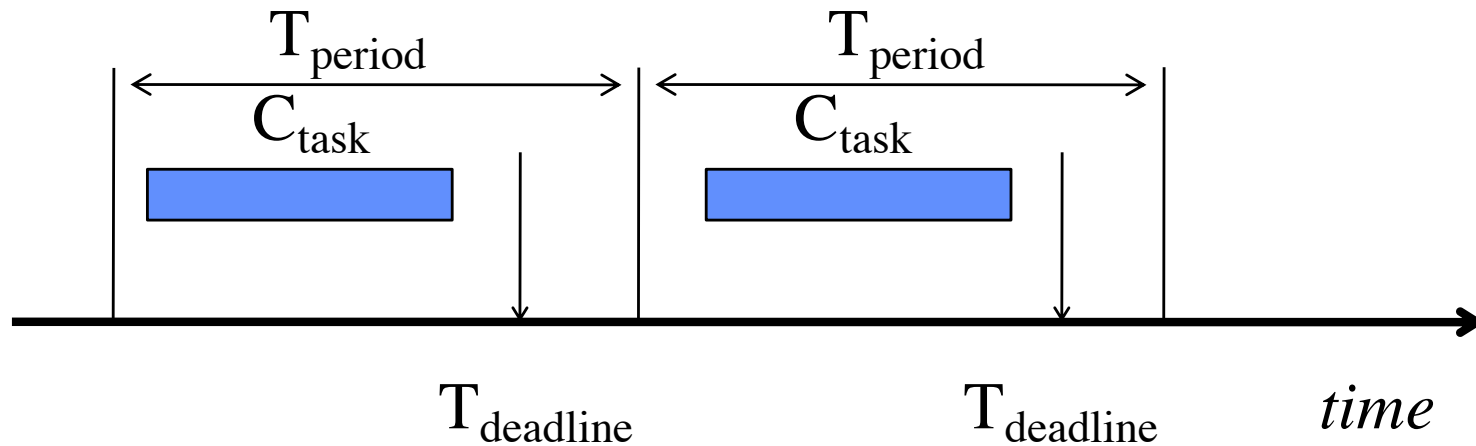
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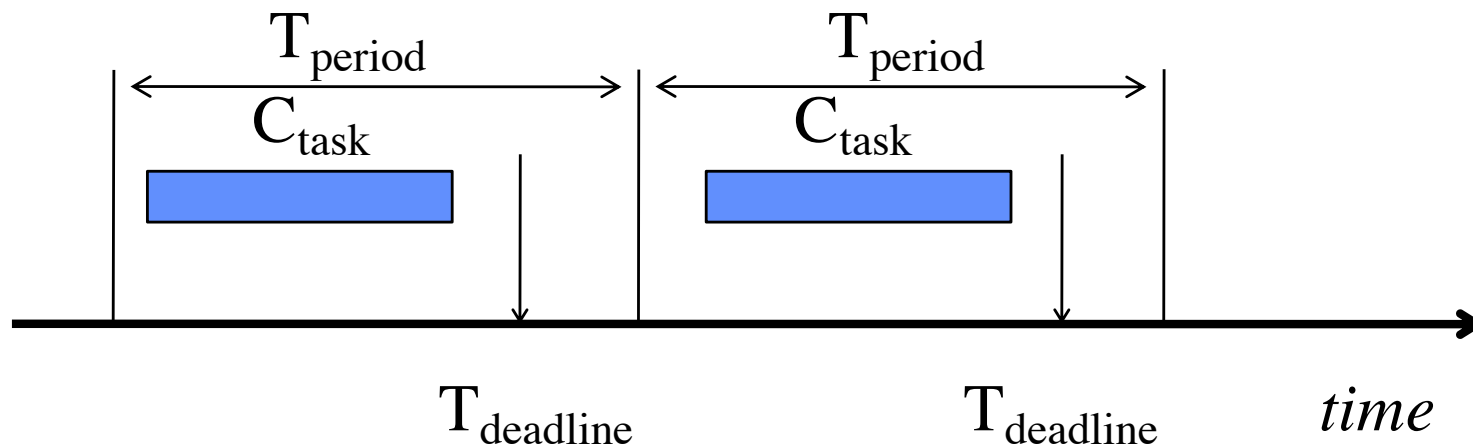
C_{task} : A Task's worst case execution time



A little theory....



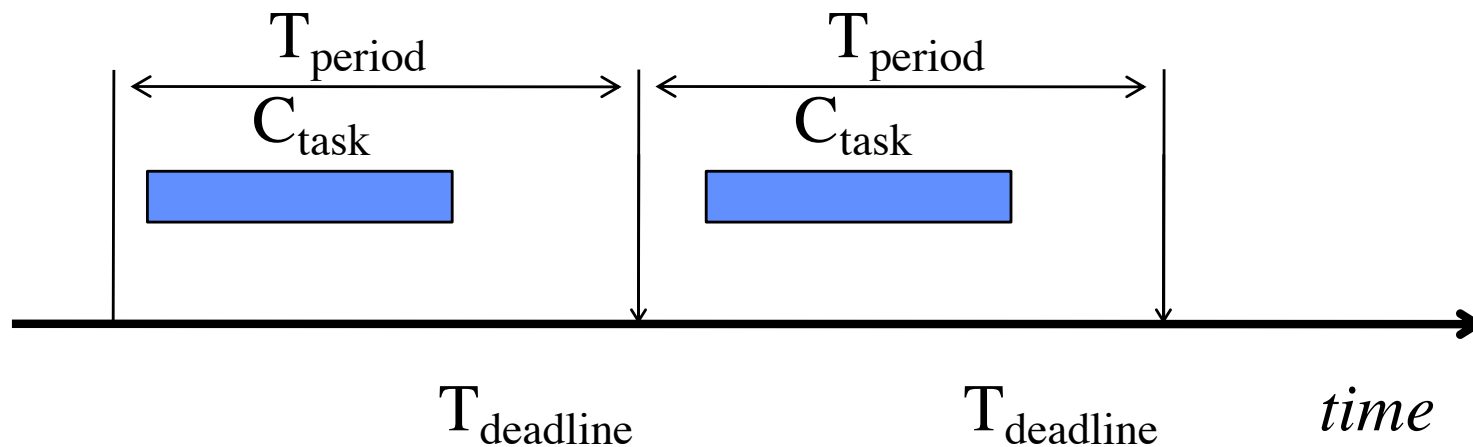
A little theory....



Task must complete by it's deadline T_{deadline}



A little theory....



Task must complete by it's deadline T_{deadline}

Simplify to make $T_{\text{deadline}} = \text{end of period}$



RT- Scheduling Algorithms

- Priority based scheduling
 - static priority;
 - Priority set during design time
 - Does not change during system operation
 - Dynamic priorities
 - Change as system runs
- Preemption/Non-Preemption
 - Preemption: Task on CPU can get booted by higher Priority Task ready to run
 - Non-preemption: Task on CPUs keeping executing even if higher priority task ready to run



Priority-driven scheduling example

- **Rules:**
 - each process has a fixed priority (1 highest);
 - highest-priority ready process gets CPU;
 - process continues until done or wait state.
- **Processes**
 - P1: priority 1, execution time 10
 - P2: priority 2, execution time 30
 - P3: priority 3, execution time 20



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- Rules:

- each process has a fixed priority (1 highest);
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- P1: priority 1, execution time 10 Ready at T_{15}
- P2: priority 2, execution time 30 Ready at T_0
- P3: priority 3, execution time 20



Priority-driven scheduling example

- Rules:

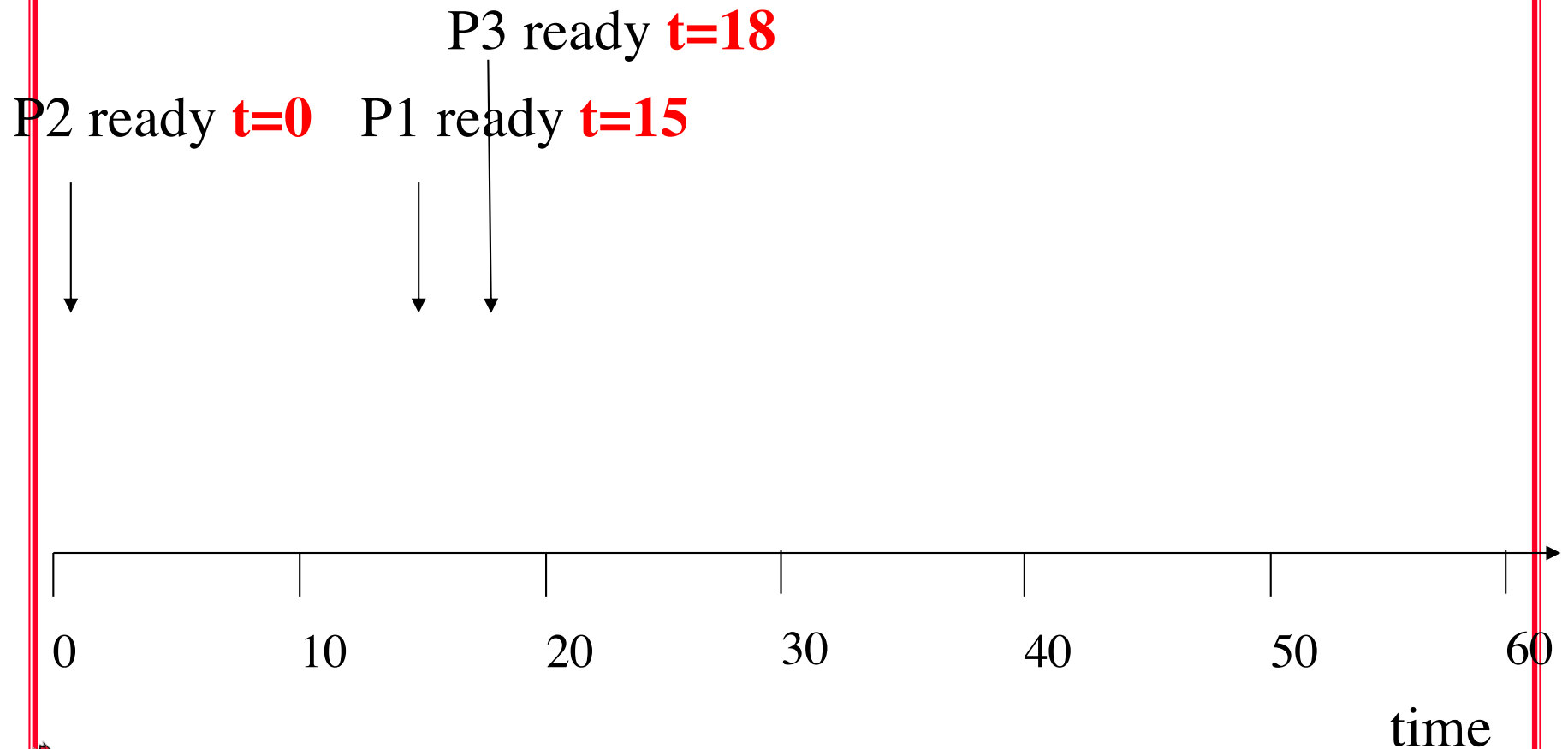
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Priority-driven scheduling example

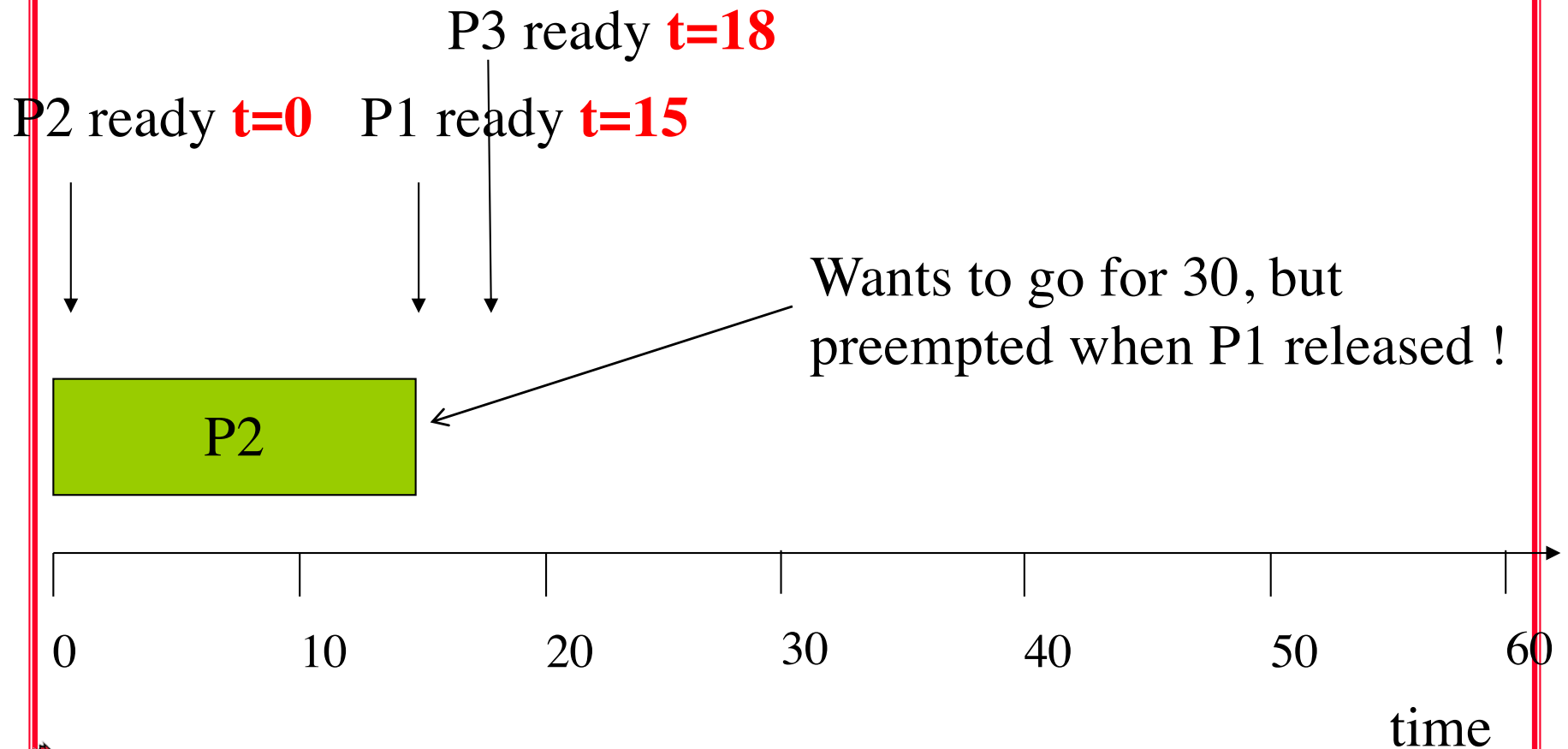


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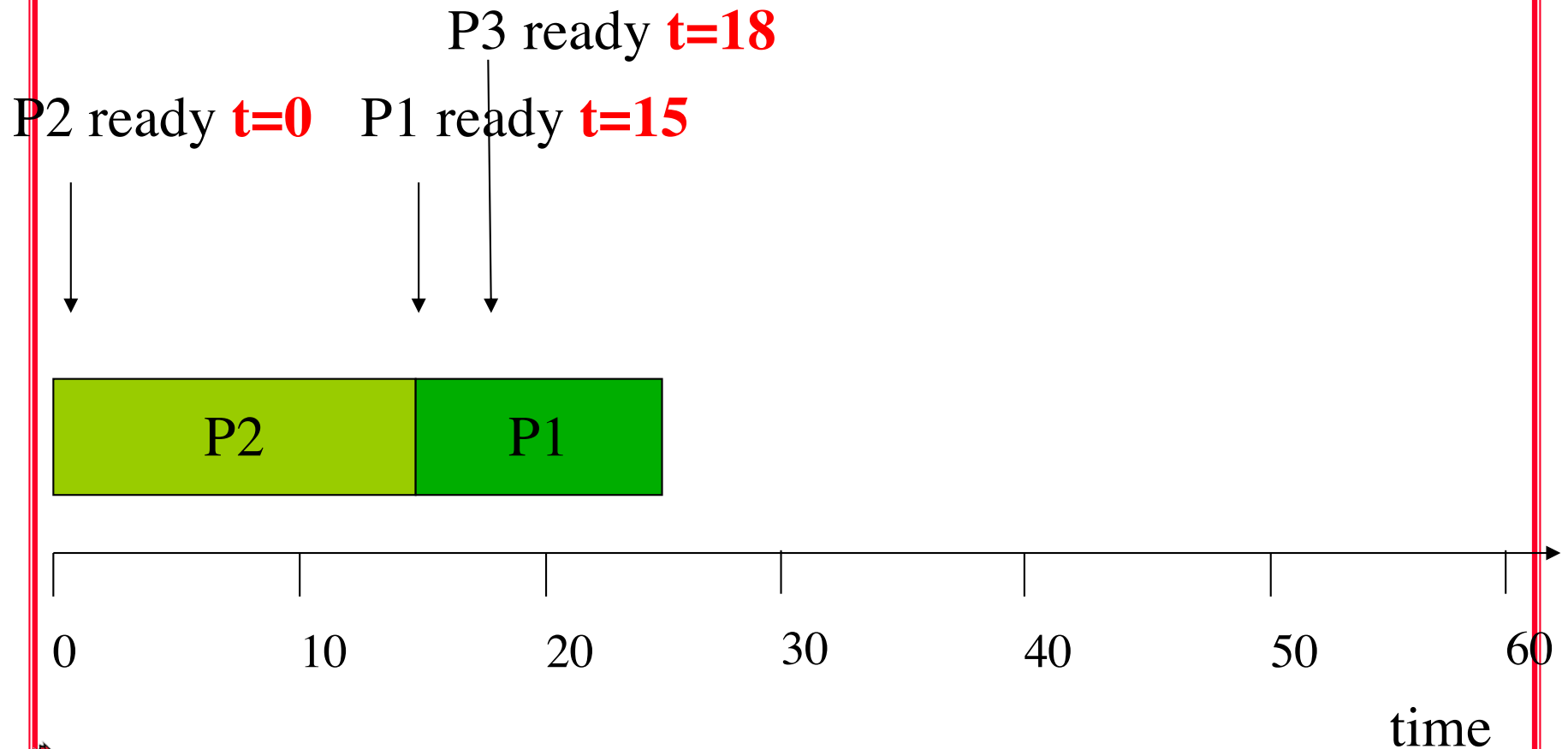
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Priority-driven scheduling example



Priority-driven scheduling example

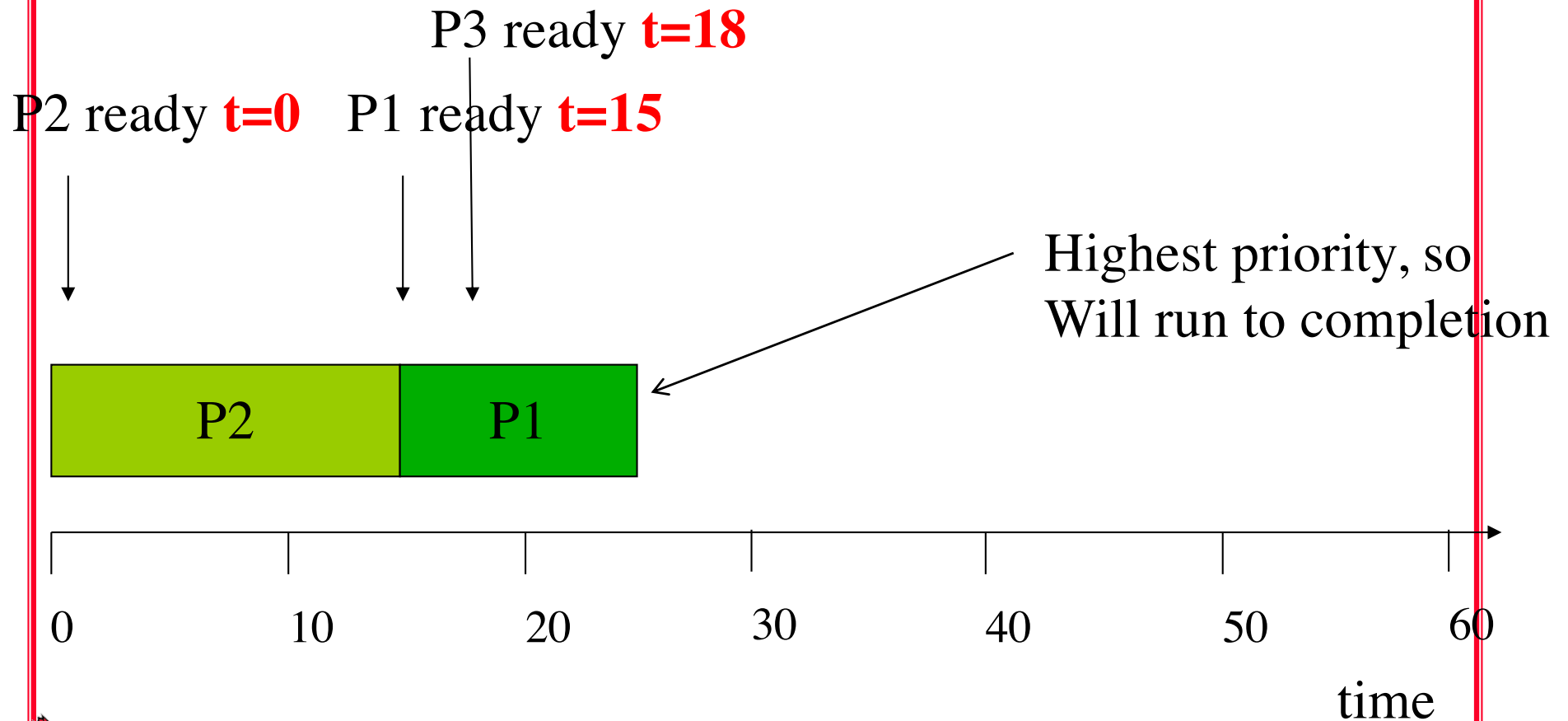


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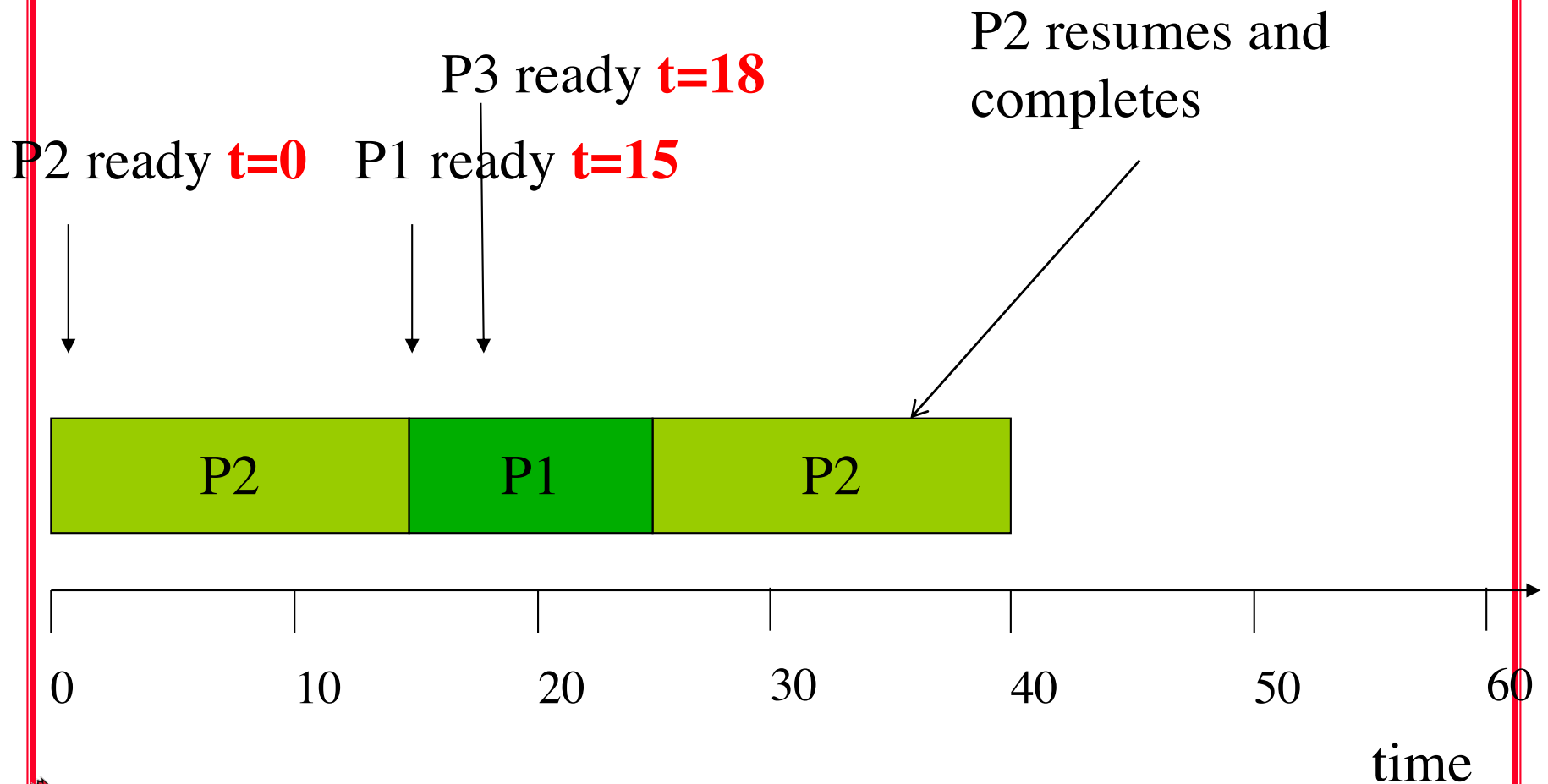


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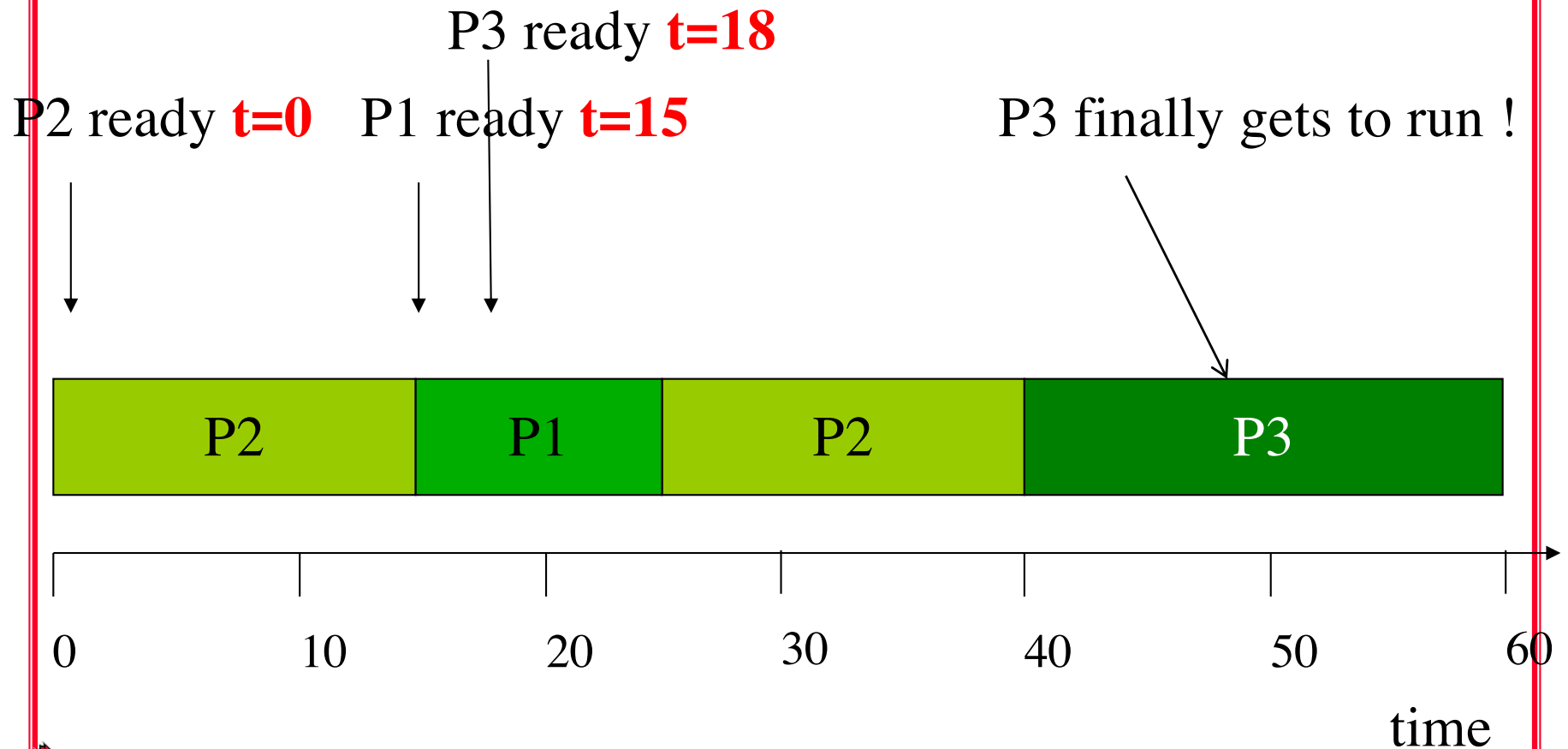
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